

Understanding Consistency Maintenance in Service Discovery Architectures in Response to Message Loss

Chris Dabrowski, Kevin Mills, Jesse Elder

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Dynamic discovery protocols in essence...

enable distributed software components

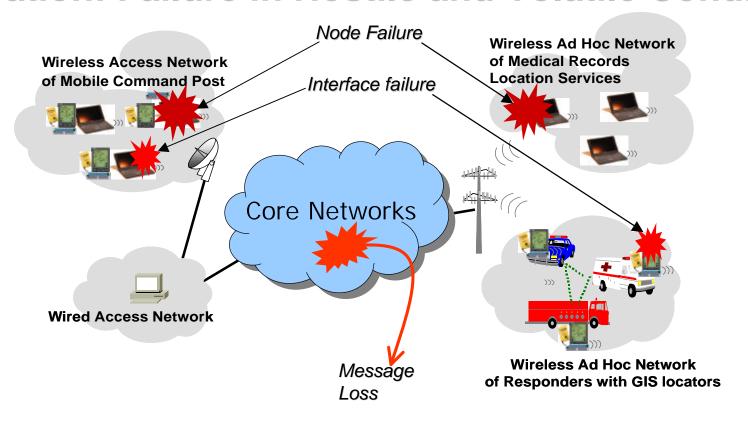
- (1) to *discover* each other without prior arrangement,
- (2) to express opportunities for collaboration,
- (3) to *compose* themselves into larger collections that cooperate to meet an application need, and
- (4) to **detect and adapt** to failures.

Some examples:





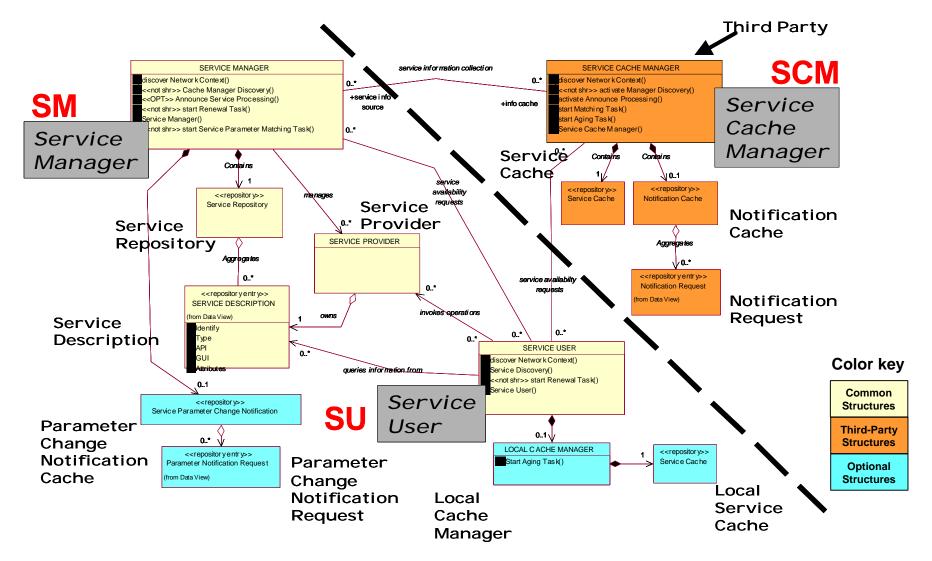
Motivation: Failure in Hostile and Volatile Conditions



 Focus of study: comparative robustness of different combinations of protocol architectures, topologies, and consistency maintenance mechanisms during message loss.

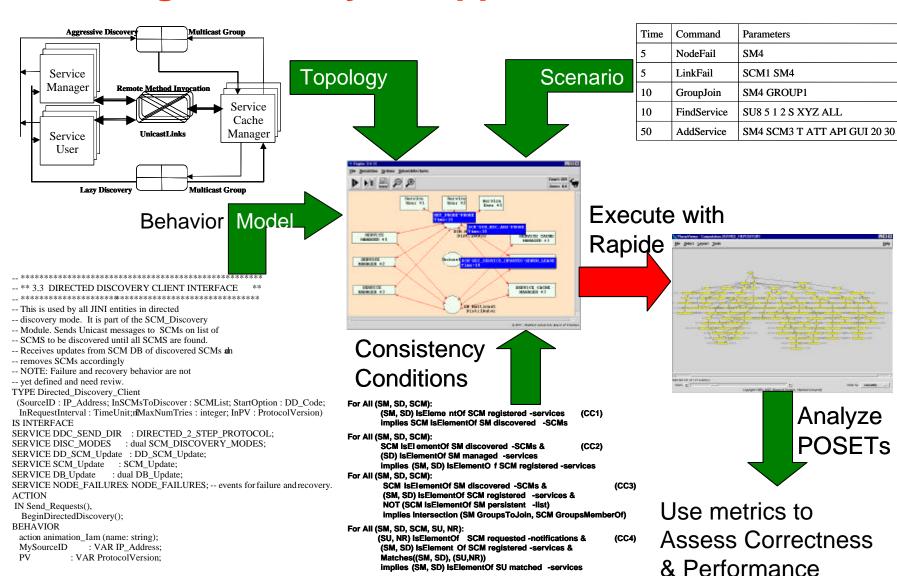


General Architecture for Service Discovery Systems



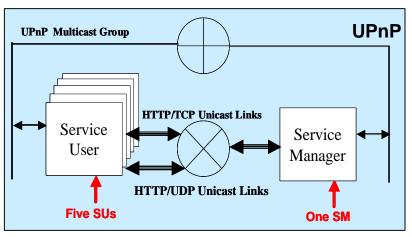


Modeling and Analysis Approach



INFORMATION TECHNOLOGY

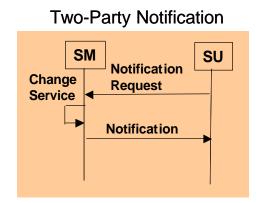
How do various service discovery architectures, topologies, and consistency-maintenance mechanisms perform under deadline during message loss conditions?



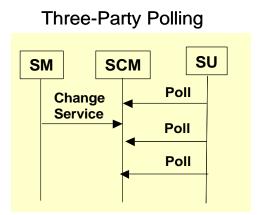
Two-Party Polling

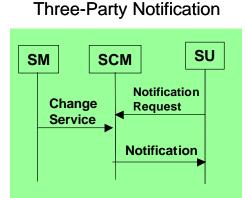
SM SU

Change Poll
Poll
Poll



Jini **Aggressive Discovery Multicast Group** One SM Optional 2nd SCM Service Manager Remote Method Invocation Service Cache Unicast Links Manager Service User One SCM Lazy Discovery Multicast Group **Five SUs**



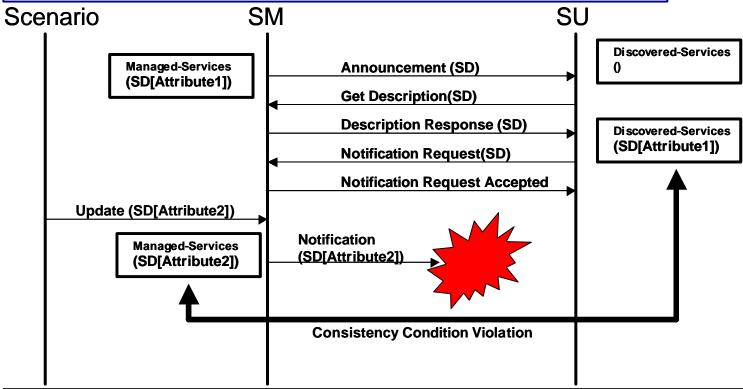




Monitoring Consistency

For All (SM, SU, SD):

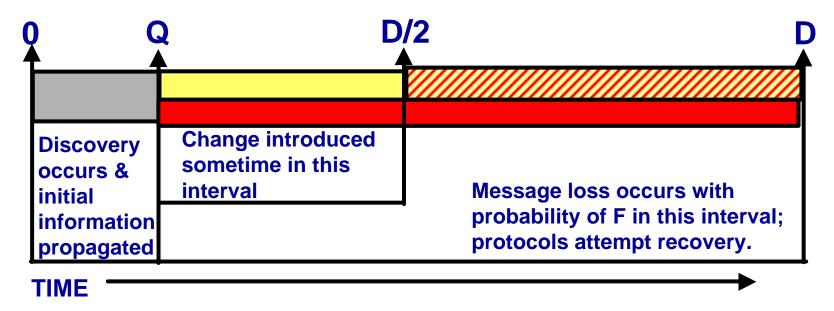
(SM, SD [Attributes1]) IsElementOf SU discovered-services
SD [Attributes2] IsElementOf SM managed-services
implies Attributes1 = Attributes2



How well does the system restore consistency after message loss?



Modeling Message Loss



Random Processes

- Choose a time to introduce the change [uniform(Q, D/2)]
- 2. For each message transmission, determine if message is lost using F

Q = end of quiescent period (100 s in our experiment)

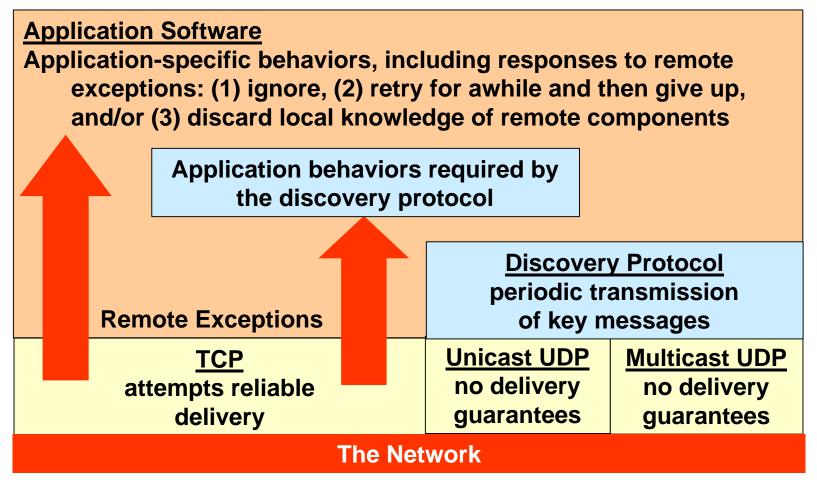
D = propagation deadline (5400 s in our experiment)

F = message loss rate (variable from 0% - 95% in 5% increments in our experiment)



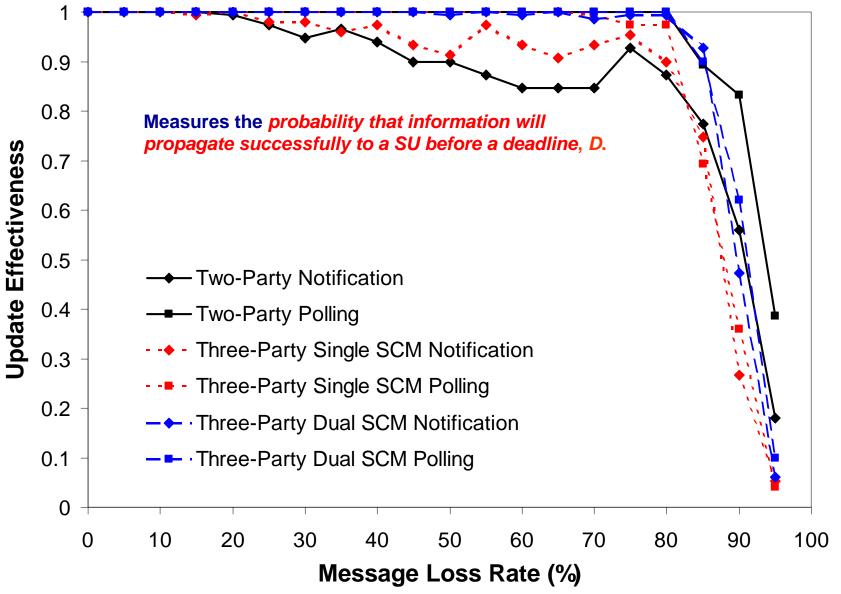
Division of Failure Recovery Responsibilities:

Communication Protocol - Discovery Protocol - Application Software



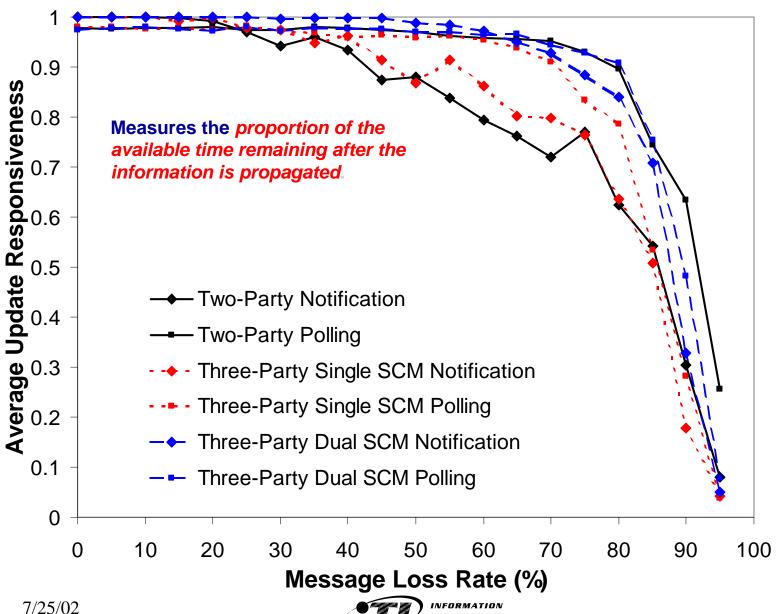


Update Effectiveness UPnP (2-Party) vs. Jini (3-Party)

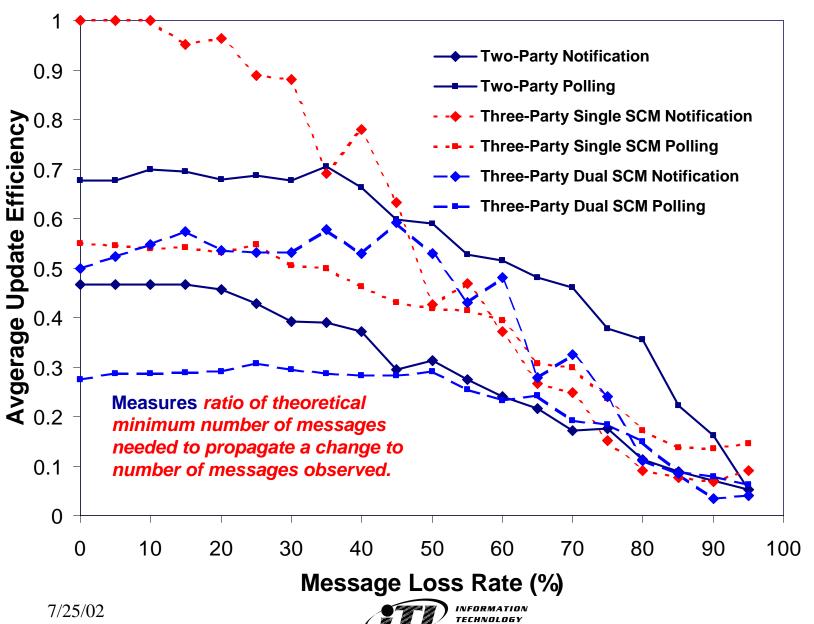




Update Responsiveness UPnP (2-Party) vs. Jini (3-Party)



Update Efficiency UPnP (2-Party) vs. Jini (3-Party)



Conclusions

- Executable architectural models represent essential complexity and reveal collective dynamics – leading to valuable insights into performance of middleware services during message loss
 - paper specifications do not represent dynamics very well
 - reference implementations exhibit substantial incidental complexity
- A single architectural model can be analyzed for behavioral, performance, and logical properties
 - limits errors and inconsistencies that can creep in when using multiple models to represent different facets of a design
- Due to effect of recovery strategies, 2- and 3-party discovery architectures exhibit similar robustness during message loss, but
 - Overall, polling provides slightly better effectiveness and responsiveness but lowers efficiency, but with significant exceptions.
 - Sole reliance on TCP retransmissions to recover notifications leads to an unexpected reduction in update effectiveness at lower message loss rates, which rises at higher rates as recovery strategies come into play (most pronounced for UPnP; Jini has additional SM behaviors that compensate).

